LEAN THEOREM PROVER

# past, present and future

Galois Inc., Oregon, August 2018

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https://leanprover.github.io

*Lean is a platform for software verification and formalized mathematics*

# Goals

- Proof stability
- Extensibility
- Expressivity - Dependent Type Theory
- Scalability
- de Bruijn’s principle: small trusted kernel, and 2 external type checkers

“Hack without fear”

# Motivation: automated provers @ Microsoft

Testing

SAGE

Pex

Software Verification

F*

Vcc

Alive Ivy

BOOGIE

Z3 Theorem Prover

# Software verification & automated provers

- Easy to use for simple properties
- Main problems:
  - Scalability issues
  - Proof stability
  - Hard to control the behavior of automated provers
- in many verification projects:
  - Hyper-V
  - Ironclad & Ironfleet (https://github.com/Microsoft/Ironclad)
  - Everest (https://project-everest.github.io/)

## *Extend Lean using Lean*

Metaprogramming

Domain specific automation

Domain specific languages

## *Whitebox automation*

Access Lean internals using Lean

Simplifiers, decision procedures, type class resolution, type inference, unifiers, matchers, …

# Applications

- IVy Metatheory (Ken McMillan - MSR Redmond)
- AliveInLean (Nuno Lopes - MSR Cambridge)
- Protocol Verification (Joe Hendrix, Joey Dodds, Ben Sherman, Ledah Casburn, Simon Hudon - Galois)
- Verified Machine Learning (Daniel Selsam - Stanford)
- SQL query equivalence (Shumo Chu et al - UW)

# Applications (cont.)

- FormalAbstracts (Tom Hales - University of Pittsburgh)
- Lean Forward, Number Theory (Jasmin Blanchette - Vrije Universiteit)
- Mathlib (Mario Carneiro - CMU and Johannes Hölzl - Vrije Universiteit)
- Teaching
    - Logic and Automated Reasoning (Jeremy Avigad - CMU)
    - Programming Languages (Zach Tatlock - UW)
    - Foundations of Analysis (Kevin Buzzard - Imperial College)

# Alive

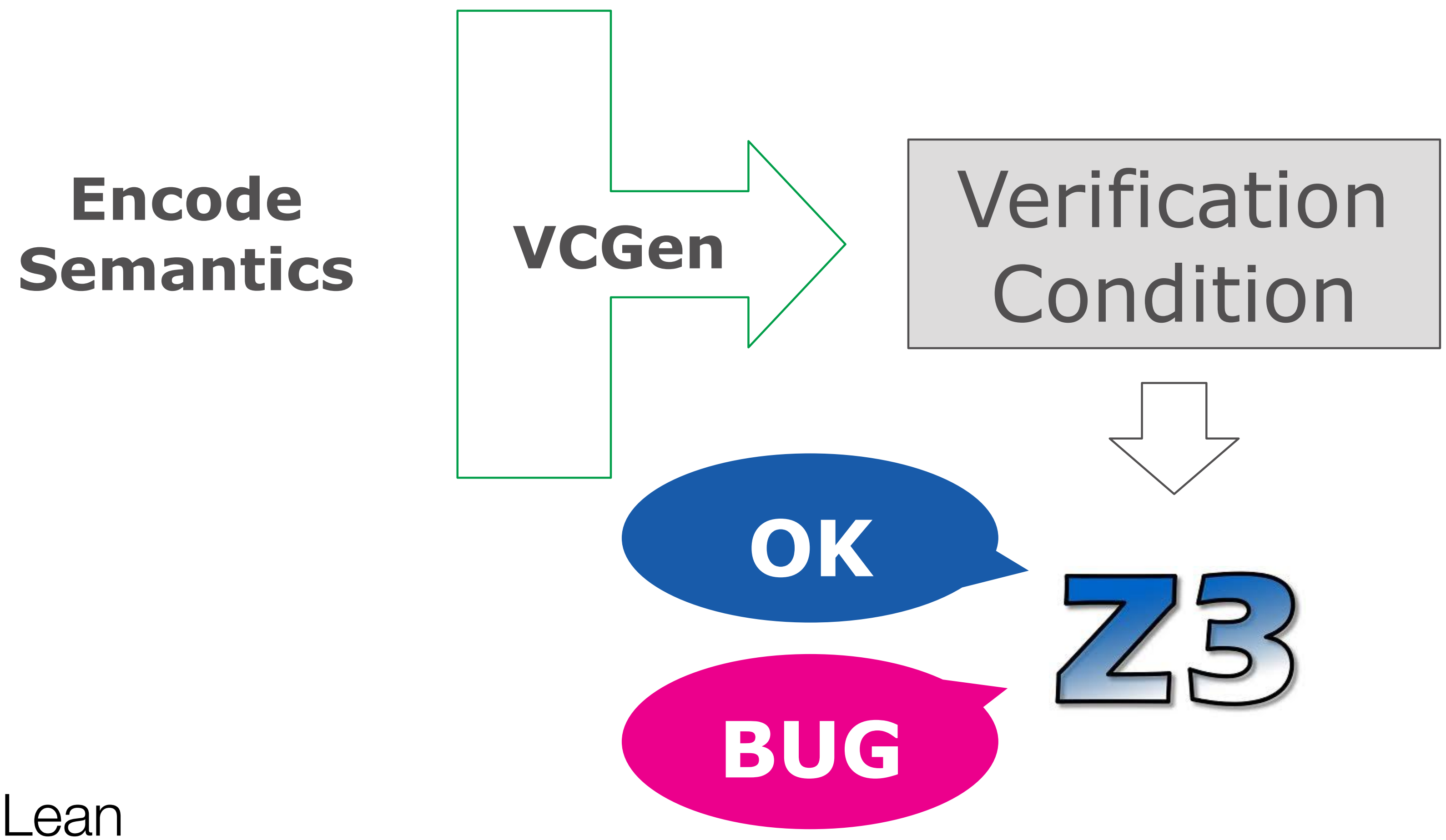
Nuno Lopes, MSR Cambridge

```
Pre: isPowerOf2(C)
%s = shl C, %N
%q = zext %s
%r = udiv %x, %q

  =>

%N2 = add %N, log2(C)
%N3 = zext %N2
%r  = lshr %x, %N3
```

<Input>




Re-implementation of Alive in Lean

Open source: https://github.com/Microsoft/AliveInLean

Pending issues:

- Using processes+pipes to communicate with Z3
- Simpler framework for specifying LLVM instructions

*Lean Demo*

Writing metaprograms/tactics/automation in Lean

# Metaprogramming example

```
meta def find : expr → list expr → tactic expr
| e []          := failed
| e (h :: hs) :=
  do t ← infer_type h,
     (unify e t >> return h) <|> find e hs

meta def assumption : tactic unit :=
do { ctx ← local_context,
     t   ← target,
     h   ← find t ctx,
     exact h }
<|> fail "assumption tactic failed"

lemma simple (p q : Prop) (h₁ : p) (h₂ : q) : q :=
by assumption
```

# Reflecting expressions

```
inductive level
| zero   : level
| succ   : level → level
| max    : level → level → level
| imax   : level → level → level
| param  : name → level
| mvar   : name → level
```

```
inductive expr
| var    : nat → expr
| lconst : name → name → expr
| mvar   : name → expr → expr
| sort   : level → expr
| const  : name → list level → expr
| app    : expr → expr → expr
| lam    : name → binfo → expr → expr → expr
| pi     : name → binfo → expr → expr → expr
| elet   : name → expr → expr → expr → expr
```

```
meta def num_args : expr → nat
| (app f a) := num_args f + 1
| e         := 0
```

# Quotations

```
example : true ∧ true :=
by do apply `(and.intro trivial trivial)
```

```
example (p : Prop) : p → p ∨ false :=
by do e ← intro `h, refine ``(or.inl %%e)
```

```
meta def is_not : expr → option expr
| `(not %%a)         := some a
| `(%%a → false)     := some a
| _                  := none
```

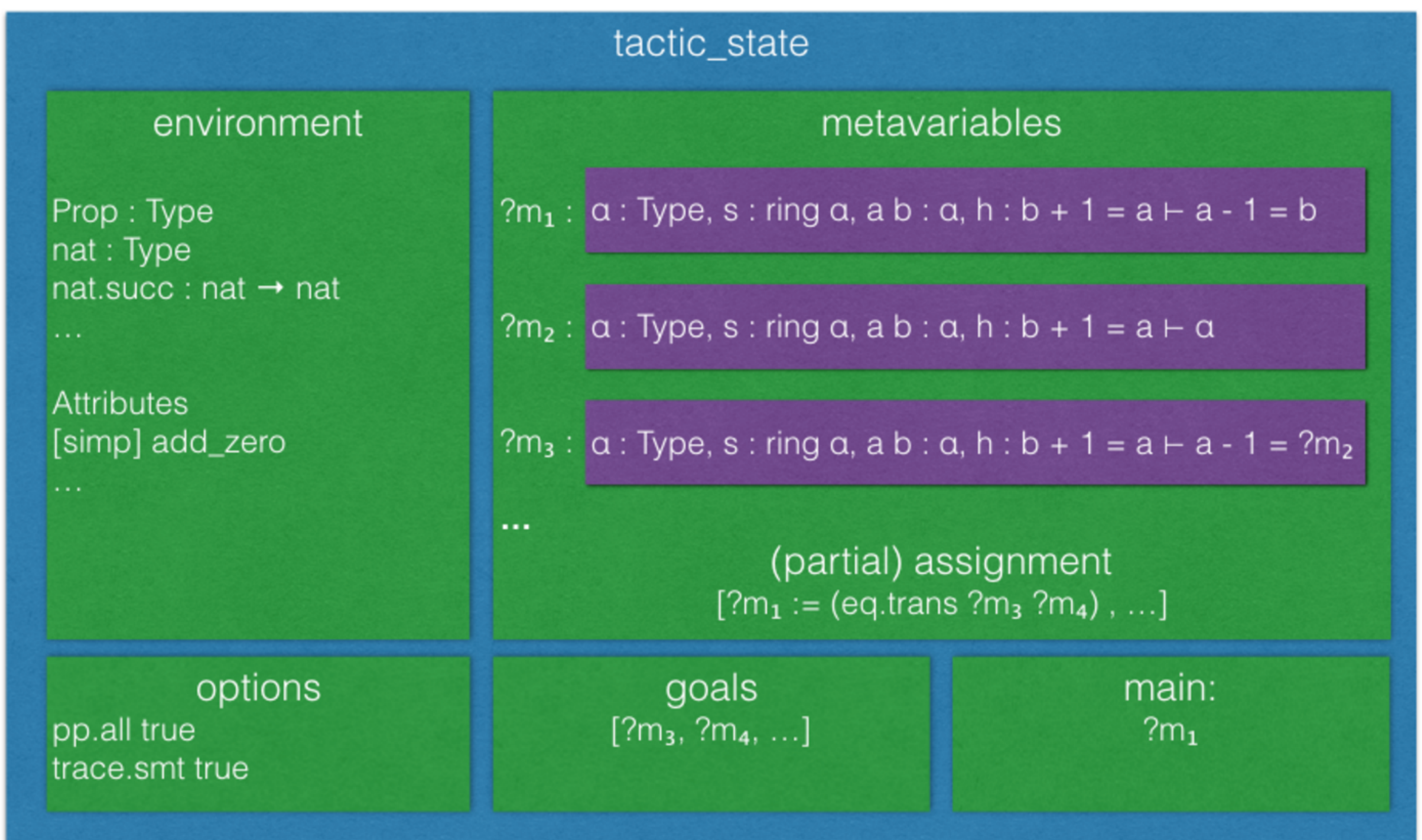
```
meta def is_not : expr → option expr
| (app (const ``not _) a)           := some a
| (pi _ _ a (const ``false _))      := some a
| _                                 := none
```

# The tactic monad

```
meta inductive result (state : Type) (α : Type)
| success   : α → state → result
| exception : option (unit → format) → option pos → state → result

meta def interaction_monad (state : Type) (α : Type) :=
state → result state α

meta def tactic := interaction_monad tactic_state
```



# Extending the tactic state

```
def state_t (σ : Type) (m : Type → Type) [monad m] (α : Type) : Type :=
σ → m (α × σ)

meta constant smt_goal : Type
meta def smt_state  := list smt_goal
meta def smt_tactic := state_t smt_state tactic

meta def eblast : smt_tactic unit := repeat (ematch; try close)

meta def collect_implied_eqs : tactic cc_state :=
focus $ using_smt $ do
  add_lemmas_from_facts, eblast,
  (done; return cc_state.mk) <|> to_cc_state
```

# Superposition prover

- 2200 lines of code

```
example {α} [monoid α] [has_inv α] : (∀ x : α, x * x⁻¹ = 1) →
                                      ∀ x : α, x⁻¹ * x = 1 :=
by super with mul_assoc mul_one

meta structure prover_state :=
(active passive : rb_map clause_id derived_clause)
(newly_derived : list derived_clause) (prec : list expr)
(locked : list locked_clause) (sat_solver : cdcl.state)
...
meta def prover := state_t prover_state tactic
```

# dlist

```
structure dlist (α : Type u) :=
(apply     : list α → list α)
(invariant : ∀ l, apply l = apply [] ++ l)
```

```
def to_list : dlist α → list α
| ⟨xs, _⟩ := xs []
```

```
local notation `♯`:max := by abstract {intros, rsimp}
```

```
/-- `O(1)` Append dlists -/
protected def append : dlist α → dlist α → dlist α
| ⟨xs, h₁⟩ ⟨ys, h₂⟩ := ⟨xs ∘ ys, ♯⟩

instance : has_append (dlist α) :=
⟨dlist.append⟩
```

# transfer tactic

- Developed by Johannes Hölzl (approx. 200 lines of code)

```
lemma to_list_append (l₁ l₂ : dlist α) : to_list (l₁ ++ l₂) = to_list l₁ ++ to_list l₂ :=
show to_list (dlist.append l₁ l₂) = to_list l₁ ++ to_list l₂, from
by cases l₁; cases l₂; simp; rsimp
```

```
protected def rel_dlist_list (d : dlist α) (l : list α) : Prop :=
to_list d = l
```

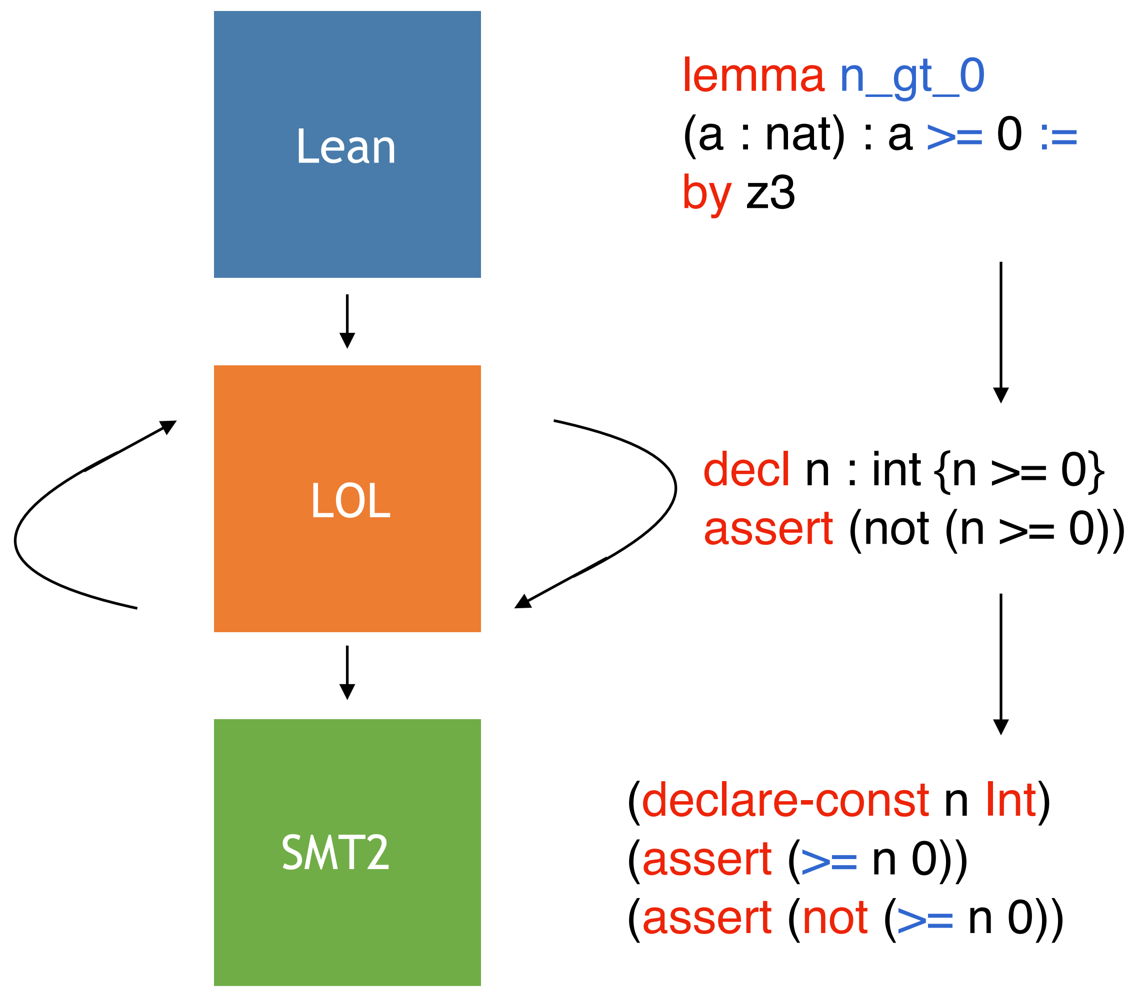
```
protected meta def transfer : tactic unit := do
  _root_.transfer.transfer [`relator.rel_forall_of_total, `dlist.rel_eq, `dlist.rel_empty,
    `dlist.rel_singleton, `dlist.rel_append, `dlist.rel_cons, `dlist.rel_concat]

example : ∀(a b c : dlist α), a ++ (b ++ c) = (a ++ b) ++ c :=
begin
  dlist.transfer,
  intros,
  simp
end
```

- We also use it to transfer results from nat to int.

# Lean to SMT2

- Goal: translate a Lean local context, and goal into SMT2 query.
- Recognize fragment and translate to low-order logic (LOL).
- Logic supports some higher order features, is successively lowered to FOL, finally SMT2.



```
mutual inductive type, term
with type : Type
| bool : type
| int : type
| var : string → type
| fn : list type → type → type
| refinement : type → (string → term) → type
with term : Type
| apply : string → list term → term
| true : term
| false : term
| var : string → term
| equals : term → term → term
| …
| forallq : string → type → term → term
```

```
meta structure context :=
(type_decl : rb_map string type)
(decls : rb_map string decl)
(assertions : list term)
```

```
meta def reflect_prop_formula' : expr → smt2_m lol.term
| `(¬ %%P) := lol.term.not <$> (reflect_prop_formula' P)
| `(%%P = %%Q) := lol.term.equals <$>
                (reflect_prop_formula' P) <*>
                (reflect_prop_formula' Q)
| `(%%P ∧ %%Q) := lol.term.and <$>
                (reflect_prop_formula' P) <*>
                (reflect_prop_formula' Q)
| `(%%P ∨ %%Q) := lol.term.or <$>
                (reflect_prop_formula' P) <*>
                (reflect_prop_formula' Q)
| `(%%P < %%Q) := reflect_ordering lol.term.lt P Q
| …
| `(true) := return $ lol.term.true
| `(false) := return $ lol.term.false
| e := …
```

# Coinductive predicates

- Developed by Johannes Hölzl (approx. 800 lines of code)
- Uses impredicativity of Prop
- No kernel extension is needed

```
coinductive all_stream {α : Type u} (s : set α) : stream α → Prop
| step {} : ∀{a : α} {ω : stream α}, a ∈ s → all_stream ω → all_stream (a :: ω)
```

```
coinductive alt_stream : stream bool → Prop
| tt_step : ∀{ω : stream bool}, alt_stream (ff :: ω) → alt_stream (tt :: ff :: ω)
| ff_step : ∀{ω : stream bool}, alt_stream (tt :: ω) → alt_stream (ff :: tt :: ω)
```

# Ring solver

- Developed by Mario Carneiro (approx. 500 lines of code)
- https://github.com/leanprover/mathlib/blob/master/tactic/ring.lean
- ring2 uses computational reflection

```
import tactic.ring

theorem ex1 (a b c d : int) : (a + 0 + b) * (c + d) = b*d + c*b + a * c + d * a :=
by ring

theorem ex2 (α : Type) [comm_ring α] (a b c d : α)
            : (a + 0 + b) * (c + d) = b*d + c*b + a * c + d * a :=
by ring
```

# Fourier-Motzkin elimination

- Linear arithmetic inequalities
- Developed here
- https://github.com/GaloisInc/lean-protocol-support/tree/master/galois/arith

# Lean 3.x limitations

- Lean programs are compiled into byte code
- Lean expressions are foreign objects in the Lean VM
- Very limited ways to extend the parser

```
infix >=         := ge
infix ≥          := ge
infix >          := gt
```

```
notation `∃` binders `, ` r:(scoped P, Exists P) := r
```

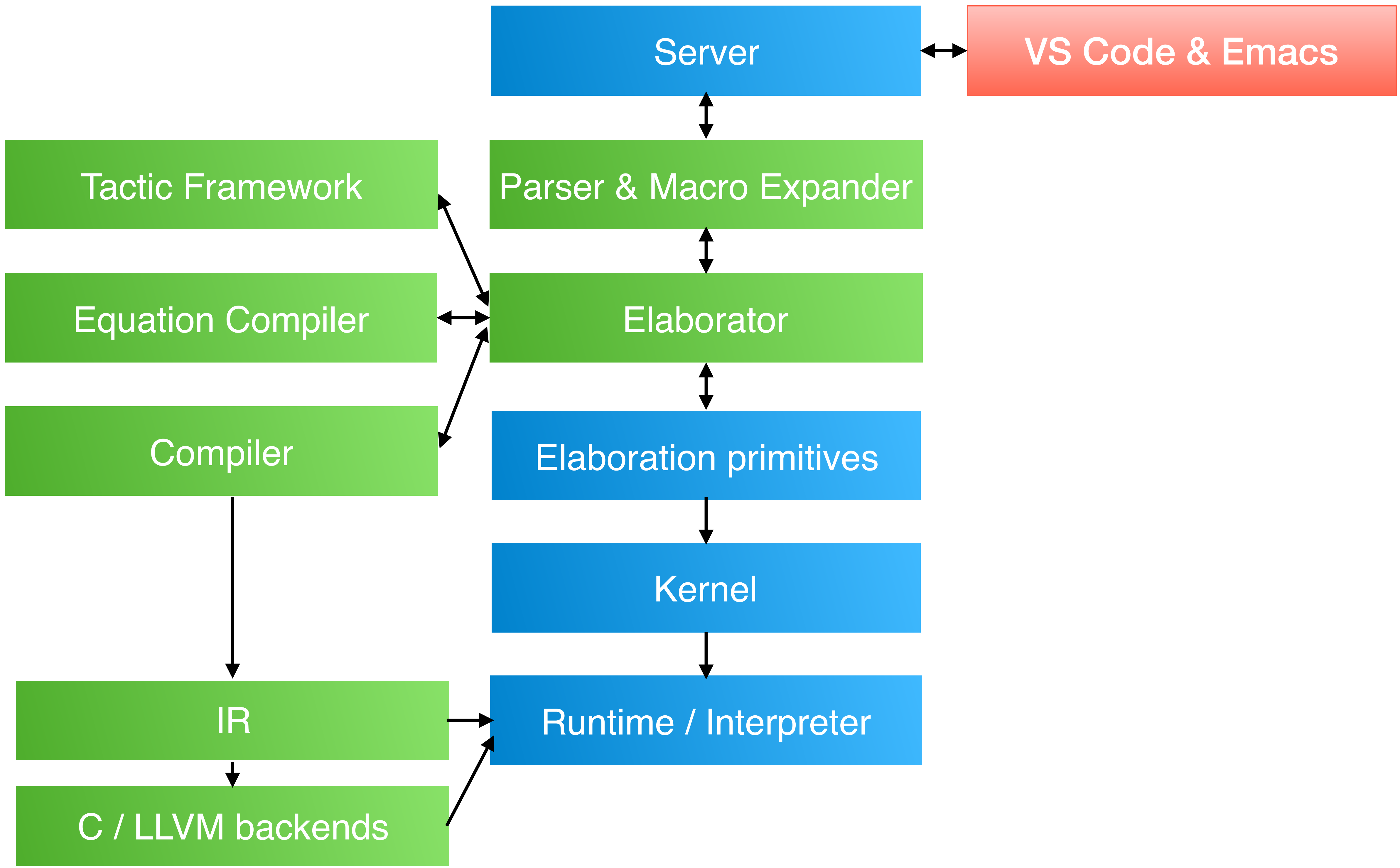
```
notation `[` l:(foldr `, ` (h t, list.cons h t) list.nil `]`) := l
```

- Users cannot implement their own elaboration strategies
- Users cannot extend the equation compiler (e.g., support for quotient types)

# Lean 4

- Leo and Sebastian Ullrich (and soon Gabriel Ebner)
- Implement Lean in Lean
  - parser, elaborator, equation compiler, code generator, tactic framework and formatter
- New intermediate representation (defined in Lean) can be translated into C++ (and LLVM IR)
- Only runtime, kernel and basic primitives are implemented in C++
- Users may want to try to prove parts of the Lean code generator or implement their own kernel in Lean
- Foreign function interface (invoke external tools)

# Lean 4 architecture



# Parser

- Implemented in Lean
- Fully extensible
- Design your own domain specific language
- Error recovery, documentation, printer, …  for free

```
@[irreducible, derive monad alternative monad_reader monad_state monad_parsec monad_except]
def read_m := rec_t syntax $ reader_t reader_config $ state_t reader_state $ parsec syntax

structure reader :=
(read : read_m syntax)
(tokens : list token_config := [])
```

```
def open_export.reader : reader :=
[ident,
 ["as", ident]?,
 [try ["(", ident], ident*, ")"]?,
 [try ["(", "renaming"], [ident, "->", ident]+, ")"]?,
 ["(", "hiding", ident+, ")"]?
]+

def open.reader : reader :=
node «open» ["open", open_export.reader]
```

# Syntax Objects

```
structure syntax_ident :=
(info : option source_info) (name : name) (msc : option macro_scope_id) (res : option resolved)

inductive atomic_val
| string (s : string)
| name   (n : name)

structure syntax_atom :=
(info : option source_info) (val : atomic_val)

structure syntax_node (syntax : Type) :=
(macro : name) (args : list syntax)

inductive syntax
| ident (val : syntax_ident)
/- any non-ident atom -/
| atom (val : syntax_atom)
| node (val : syntax_node syntax)
```

Macros can be expanded and/or elaborated.
Users can define new readers and macros.

# Kernel expressions

**Elaborator converts syntax objects into expressions.**

```
inductive expr
| bvar  : nat → expr                                         -- bound variables
| fvar  : name → expr                                        -- free variables
| mvar  : name → expr → expr                                 -- (temporary) meta variables
| sort  : level → expr                                       -- Sort
| const : name → list level → expr                           -- constants
| app   : expr → expr → expr                                 -- application
| lam   : name → binder_info → expr → expr → expr            -- lambda abstraction
| pi    : name → binder_info → expr → expr → expr            -- Pi
| elet  : name → expr → expr → expr → expr                   -- let expressions
| lit   : literal → expr                                     -- literals
| mdata : kvmap → expr → expr                                -- metadata
| proj  : nat → expr → expr                                  -- projection
```

# Compiler - code generator

- Implemented Lean
- External contributors can prove the new compiler is correct
- Code specialization and monomorphization
- Target is the new IR also defined in Lean
- Users can select theorems as optimization rules

```
@[simp] lemma map_map (g : β → γ) (f : α → β) (l : list α) : map g (map f l) = map (g ∘ f) l :=
by induction l; simp [*]
```

# Runtime

Strict, GC based on reference counting, destructive updates for unshared objects, support for unboxed values.

```
/- IR Instructions -/
inductive instr
| assign       (x : var) (ty : type) (y : var)                        -- x : ty := y
| assign_lit   (x : var) (ty : type) (lit : literal)                  -- x : ty := lit
| assign_unop  (x : var) (ty : type) (op : assign_unop) (y : var)     -- x : ty := op y
| assign_binop (x : var) (ty : type) (op : assign_binop) (y z : var)  -- x : ty := op y z
| unop         (op : unop) (x : var)                                  -- op x
| call         (xs : list var) (f : fnid) (ys : list var)             -- Function call:  xs := f ys
/- Constructor objects -/
| cnstr   (o : var) (tag : tag) (nobjs : uint16) (ssz : usize)        -- Create constructor object
| set     (o : var) (i : uint16) (x : var)                            -- Set object field:        set o i x
| get     (x : var) (o : var) (i : uint16)                            -- Get object field:        x := get o i
| sset    (o : var) (d : usize) (v : var)                             -- Set scalar field:        sset o d v
| sget    (x : var) (ty : type) (o : var) (d : usize)                 -- Get scalar field:        x : ty := sget o d
/- Closures -/
| closure (x : var) (f : fnid) (ys : list var)                        -- Create closure:          x := closure f ys
| apply   (x : var) (ys : list var)                                   -- Apply closure:           x := apply ys
/- Arrays -/
| array   (a sz c : var)                                              -- Create array of objects with size `sz` and capacity `c`
| sarray  (a : var) (ty : type) (sz c : var)                          -- Create scalar array
| array_write (a i v : var)                                           -- (scalar) Array write     write a i v
```

```
inductive unop
| inc_ref | dec_ref | dec_sref | inc | dec
| free | dealloc
| array_pop | sarray_pop
```

```
inductive assign_unop
| not | neg | ineg | nat2int | is_scalar | is_shared | is_null | cast | box | unbox
| array_copy | sarray_copy | array_size | sarray_size | string_len
| succ | tag | tag_ref
```

# Code generation hints

- Support for low-level tricks used in SMT and ATP.
  Example: pointer equality

```
def use_ptr_eq {α : Type u} {a b : α}
               (c : unit -> {r : bool // a = b → r = tt})
               : {r : bool // a = b → r = tt} :=
c ()
```

Given `@use_ptr_eq _ a b c`, compiler generates

```
if (addr_of(a) == addr_of(b)) return true;
else return c();
```

# Structured trace messages

- Why did my tactic/solver fail?
- Lean 3 has support for trace messages, but they are just a bunch of strings.
- Lean 4 will provide structured trace messages and APIs for browsing them.
- Traces will be generated on demand (improved discoverability).

```
inductive trace
| mk (msg : message) (subtraces : list trace)

def trace_map := rbmap pos trace (<)

structure trace_state :=
(opts : options)
(roots : trace_map)
(cur_pos : option pos)
(cur_traces : list trace)

def trace_t (m : Type → Type u) := state_t trace_state m

class monad_tracer (m : Type → Type u) :=
(trace_root {α} : pos → name → message → thunk (m α) → m α)
(trace_ctx {α} : name → message → thunk (m α) → m α)
```

# Better support for proofs by reflection

- Define an inductive datatype (`form`) that captures a class of formulas.
- Implement a decision procedure `dec_proc` for this class.
- Prove: `∀ (s : form) ctx, dec_proc s = tt → denote s ctx`
- The type checker has to reduce (`dec_proc s`). This is too inefficient in Lean 3.
- In Lean 4, we allow users to use the compiler + IR interpreter to reduce (`dec_proc s`).
- We still need to use the symbolic reduction engine to show that the current goal and (`denote s ctx`) are definitionally equal.
- Disadvantages: increases the size of the TCB, external type checkers will probably timeout in proofs using this feature.

*New application scenarios*

# Automated reasoning framework

- Many users use Python + SMT solver to developing automated reasoning engines (e.g., Alive is implemented in Z3Py).
- Lean 3 interpreter is already faster than Python.
- FFI in Lean 4 will provide (efficient) access to external SAT & SMT solvers and ATP.
- Many goodies not available in the Python + SMT framework:
  - Simplifiers.
  - Efficient symbolic simulation.
  - Custom automation.
  - Parsing framework + integration with IDEs (VS Code, Emacs).

# Domain Specific Languages

- Users can define and reason about their DSLs.
- Code reuse:
    - Compiler infrastructure.
    - Parsing framework.
    - Elaborator.
    - IDE integration.

# Lean as a general purpose programming language

- Lean is an extensible system: parser, elaborator, compiler, etc.
- User certified optimizations as conditional rewriting rules.
- New backends for the Lean 4 IR can be implemented in Lean.
- Foreign function interface.
- leanpkg - package management tool implemented in Lean.

# Conclusion

- Users can create their on automation, extend and customize Lean
- Domain specific automation
- Internal data structures and procedures are exposed to users
- Whitebox automation
- Lean 4 automation written in Lean will be much more efficient
- Lean 4 will be more extensible
- New application domains
  - Lean 4 as a more powerful Z3Py
  - Lean 4 as a platform for developing domain specific languages